Wireless Sensor Network based Shooter Localization

Miklos Maroti, Akos Ledeczi, Gyula Simon, Gyorgy Balogh, Branislav Kusy, Andras Nadas, Gabor Pap, Janos Sallai ISIS - Vanderbilt University

Overview

- CONOPS Support: Fast and accurate enemy shooter localization is key in reducing friendly casualties and neutralizing enemy combatants
- Different approach compared to other acoustic shooter localization systems (e.g., BulletEars, Pilar, Ferret, Boomerang)
- Ad-hoc wireless network of cheap acoustic sensors is used to accurately locate enemy shooters in urban environment
- Challenges
 - Severely resource constrained nodes and communication bandwidth
 - Extreme multipath effects in urban environment
 - Simultaneous shot resolution
- Performance
 - Average accuracy: 1 meter
 - Latency: 2 seconds
 - Simultaneous shot resolution: 6 shots per second

Technical Approach

- Detect TOA of acoustic shockwave and muzzle blast
- MICA2 mote (UC Berkeley and Crossbow)
- Acoustic sensor board (Vanderbilt):
 - Three acoustic channels (only one is used)
 - High-speed AD converters
 - FPGA for signal processing: shockwave and muzzle blast detection on board
- Timestamp of shockwave and/or muzzle blast is sent to mote
- Motes route time of arrival data to base station
- Base station fuses data, estimates shooter position and displays result

Technical Approach: System Architecture



Acoustic Sensor Board: False Alarm Rejection







1ms/div

Acoustic Sensor Board: Different Weapon Types





100µs/div

Acoustic Sensor Board

- Xilinx Spartan-II FPGA
- Three acoustic channels
 - Only one is used in shooter localization
 - Selection of the microphone
 - Max 1 MHz sampling rate, 12-bit ADC
- ZC encoding
 - Feature extraction
 - Data compression





MICA2 hardware

- 8 MHz, 8-bit microcontroller
- 4 KB of RAM, 128 KB of ROM
- Wireless communication
 - Packet is 29 bytes data, 7 bytes header
 - 30 packets per second under no collisions
 - 30-100 feet range
- Power management
 - 2-3 days of continuously operation on two AA batteries
 - 2-3 of months when sleeping
- Various sensors and actuators boards
- Cost: \$150 from XBow

Middleware Services: Time Synchronization



- Requirement
 - sound travels one foot per millisecond
 - time synch error in the whole network should be less than 1 millisecond (less than 1 ft error)
- Algorithm
 - selected leader broadcasts its time
 - receivers maintain a table of global-local time pairs
 - receivers calculate clock offset and skew using linear regression
 - receivers rebroadcast the global time
- Performance:
 - ±1.5 µsec per hop error
 - Low overhead: one timesynch round per minute (i.e., one message per minute per mote)

Time Synchronization Primitive: Time Stamping of Radio Messages





4.5 µs maximum error

12 µs maximum error

Time Synchronization: Experimental Evaluation



1 message per 30 seconds per mote



layout and links:



first leadersecond leader

A. All motes are turned on

- B. The first leader is turned off
- C. Randomly selected motes were reset every 30 seconds
- D. Half of the motes were switched off
- E. All motes were switched back on

Middleware Services: Message Routing



- Requirement: acoustic event triggers many motes at once; all need to get their data to base station with low latency
- Approach: convergecast to root using our directed flood routing framework:
 - Ad-hoc routing
 - Automatic aggregation
 - Implicit acknowledgments
- Performance: when max distance from root is 5 hops, base station receives ~15 measurements in the first second

Message Routing: Channel Behavior

- MICA2 under no load: single mote is transmitting
 Effective region (95% delivery rate) is 0-10 feet
 - Transitional region (5-95% delivery rate) is 10-40 feet
- MICA2 under heavy load: most motes transmit
 - Effective region: 5 feet, transitional region: 5-30 feet
 - 70% of the motes in the transitional region receive messages with less than 30%

Polite (never transmits) and impolite (causes collision) motes

- Use probabilistic methods: rely on the unreliable
 - It is more probable that one of the motes with less that 30% delivery rate will receive a broadcasted message than one with a higher rate
 - Do not limit the next hop to a single node
 - Long unreliable links can route messages faster than short reliable ones

Message Routing: Flooding Policies

- Each flooding policy defines a state machine that describe the life cycle of data packets
- On each node each data packet is in one of the states
- Actions: received, sent, aged
- States are numbered from 0 (initial state) to 255 (final state)
- Packets with low numbered states are more important
- Packets with even numbered states are eligible for transmission



Fat Spanning-Tree Convergecast

- A spanning tree is formed
- Each node needs to know the node ID of its parent, grandparent, great-grandparent, and great-great-grandparent
- The "rank" of the node is the node ID of its grandparent
- If the rank of the sender is
 - the node ID of the grandparent of the receiver, then the sender is at the same distance
 - the node ID of the receiver or its parent, then the sender is further from the root than the receiver
 - the node ID of the great- or great-great-grand parent of the receiver, then the sender is closer to the root
 - non of the above: not in the same channel or further away.
- Increases the reliability and robustness of tree routing protocols







Middleware: Self Localization



Sensor Fusion: Consistency Surface



Sensor Fusion: Multiple Shots

- Find the global maximum of the consistency surface
- If the global maximum is above a threshold then this position is the 1st estimated shooter position
- Remove the corresponding measurements from the data set
- Recalculate surface/ global maximum (next shots)
- Shot and its detected consistent echo have the same shot time → echo elimination



System Performance

- Latency: 2 seconds
- Average accuracy in 2D (x,y): 0.64 meter; in 3D (x,y,z): 1.5 meter
- Results below based on: 71 single SRTA shots; 20 different positions (McKenna MOUT Site, Ft. Benning); 60-mote network; 100x40 m area





Experiment Results (SRTA)

Location error in meters

Location error in meters

Observed Hardware Failures



- Noisy or non-functional ADC (3%)
- Memory corruption (only one mote)
- One-way radio, only receive or transmit (3%)
- Battery exploded (only one, left in the programming board)
- Corrupted fuses (5%)
- Non-functional LEDs (1%)
- Motes overrun by cars and people (2%)
- Under direct sunlight the ADC and CPU clock of the motes changed their characteristics

Development Tools



Unionized NEST Workers



Middleware Development in Tennessee



Oops, we have to recruit a new graduate student again

